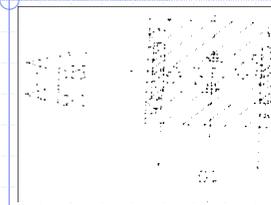


Efficient algorithms for polygonal approximation

Alexander Kolesnikov

Examples of polygonal approximation



Vectorization

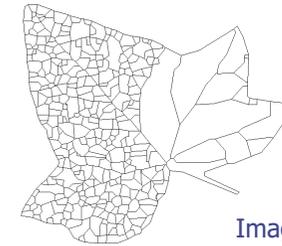
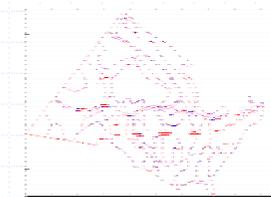
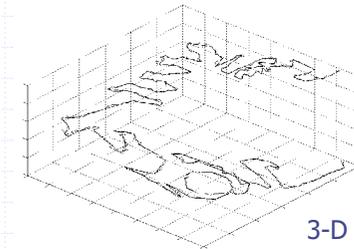


Image analysis



Digital cartography



3-D paths

Heuristic algorithms for approximation

- 1) Sequential tracing approach
- 2) Split method
- 3) Merge method
- 4) Split-and-Merge method
- 5) Dominant point detection
- 6) Relaxation labeling
- 7) K -means method
- 8) Genetic (evolutional) algorithms
- 9) Ant colony optimization method
- 10) Tabu search
- 11) Discrete particle swarm algorithm
- 12) Vertex adjustment method

Min- ϵ problem: Motivation

Heuristic algorithms

Non-optimal: $F < 100\%$

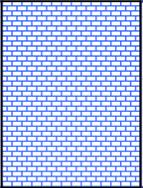
Fast: $O(N) - O(N^2)$
(seconds and less)

Optimal algorithm

Optimal: $F = 100\%$

Slow: $O(N^2) - O(N^3)$
(minutes and more)

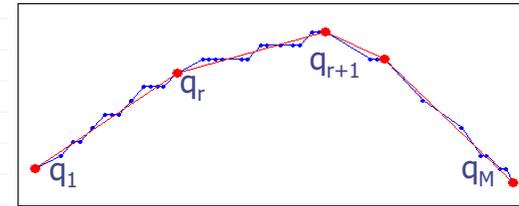
Min- ϵ problem: Motivation

Heuristic algorithms Non-optimal: $F < 100\%$ Fast: $O(N) - O(N^2)$ (seconds and less)		Optimal algorithm Optimal: $F = 100\%$ Slow: $O(N^2) - O(N^3)$ (minutes and more)
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Efficient algorithm

Fast: $O(N) - O(N^2)$
 Close to optimal: $F \approx 100\%$

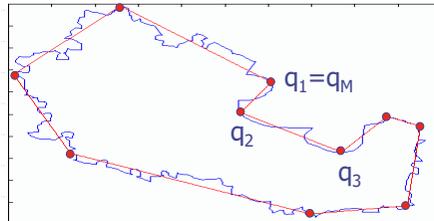
Min- ϵ problem for open curve



Approximate the given open N -vertex polygonal curve P by another one Q consisting of at most M line segments with minimum error $E(P)$:

$$E(P) = \min_{\{q_r\}} \sum_{r=2}^{M-2} e^2(q_r, q_{r+1})$$

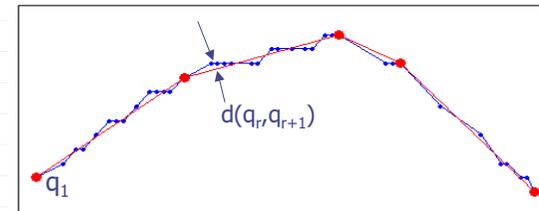
Min- ϵ problem for closed curves



Approximate the given closed N -vertex polygonal curve P by another one Q consisting of at most M line segments with minimum error $E(P)$:

$$E(P) = \min_{\{q_r\}} \sum_{r=1}^{M-1} e^2(q_r, q_{r+1})$$

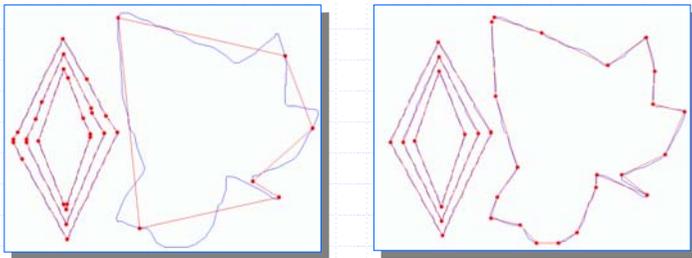
Min-# problem for open curve



Given polygonal curve P , approximate it by another polygonal curve Q with the minimum number of segments M so that the approximation error (distortion) does not exceed a given maximum tolerance ϵ :

$$M \rightarrow \min \quad \text{subject to: } d(P) \leq \epsilon$$

Multi-object $\min\text{-}\varepsilon$ problem



Given K polygonal curves P_1, P_2, \dots, P_K , approximate it by set of K other polygonal curves Q_1, Q_2, \dots, Q_K with a given total number of segments M .

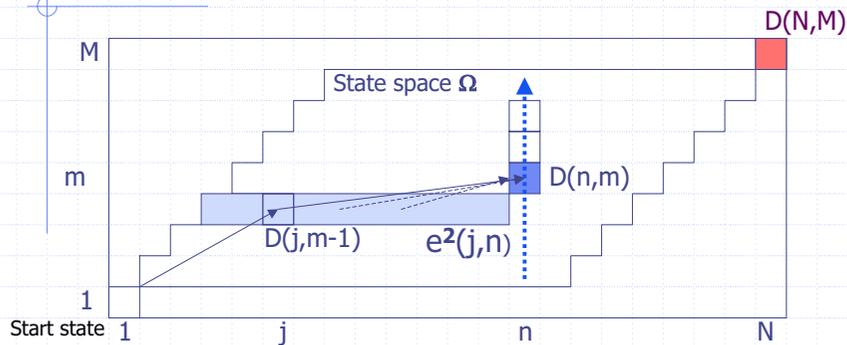
Multi-object $\min\text{-}\varepsilon$ problem (cont'd)

Given K polygonal curves P_1, P_2, \dots, P_K , approximate it by set of K other polygonal curves Q_1, Q_2, \dots, Q_K with a given total number of segments M so that the total approximation error with measure L_2 is minimized.

$$E(P_1, \dots, P_K, M) = \min_{\{M_k\}} \min_{\{q_m\}} \sum_{k=1}^K \sum_{m=1}^{M_k-1} e^2(q_{k,m}, q_{k,m+1})$$

subject to: $\sum_{k=1}^K M_k \leq M$

DP algorithm for $\min\text{-}\varepsilon$ problem



Dynamic programming algorithm:

$$D(n, m) = \min \{ D(j, m-1) + e^2(j, n) \mid m \leq j < n \}$$

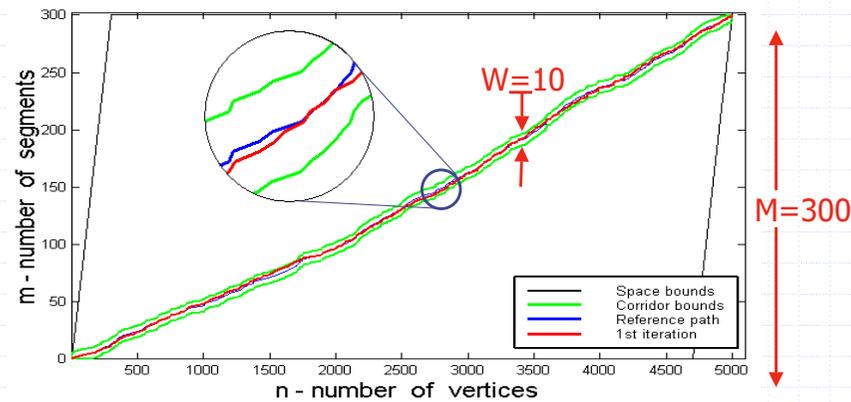
$$\forall (n, m) \in \Omega$$

Iterative reduced search for $\min\text{-}\varepsilon$ problem

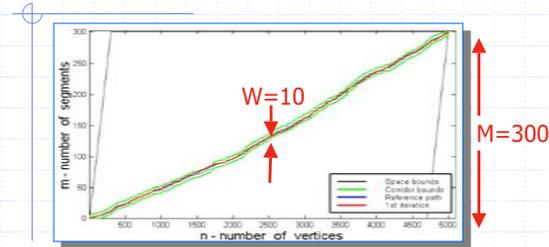
- Step 1: Find preliminary approximation with any fast heuristic algorithm.
- Step 2: Construct bounding corridor in state space along the reference path
- Step 3: Perform DP search in the bounding corridor

Use the output as a reference solution and repeat the search

Iterative reduced search DP algorithm



Iterative reduced search DP algorithm



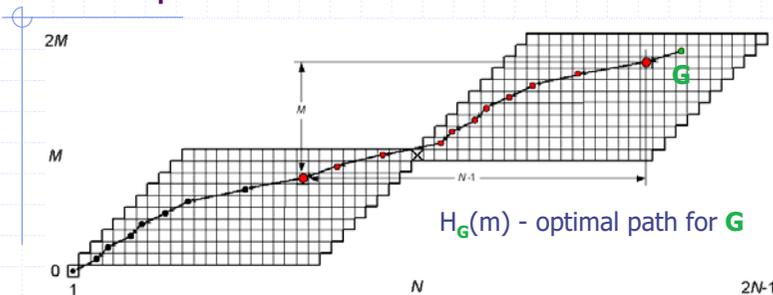
Complexity: $O(W^2N^2/M) = O(N) - O(N^2)$

Speed-up: $(W/M)^2$

Example: $(W/M)^2 = (10/300)^2 = 1/900$

Fidelity: up to 100% for iterative search

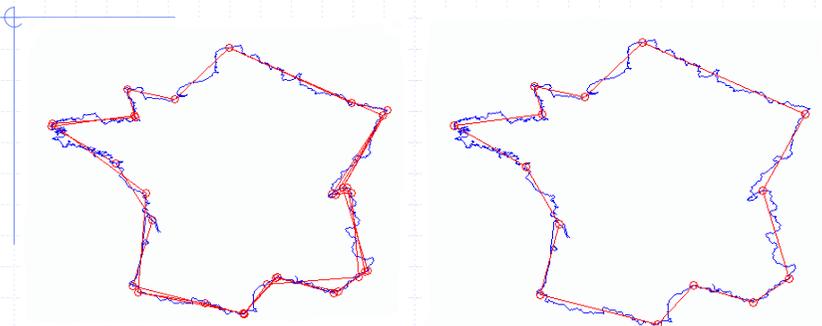
Min- ε problem for closed curves



Conjugate states: $H_G(m-M) = H_G(m) - (N-1)$;

$n_{\text{start}} = \arg \min_{M \leq m \leq 2M} \{D(H_G(m), m) - D(H_G(m-M), m-M)\} - (N-1)$;

Min- ε problem for closed curves: results



Heuristic algorithm:
Fidelity $F=88-100\%$
 $T=100$ s

Proposed algorithm:
Fidelity $F=100\%$
 $T=10.4$ s.

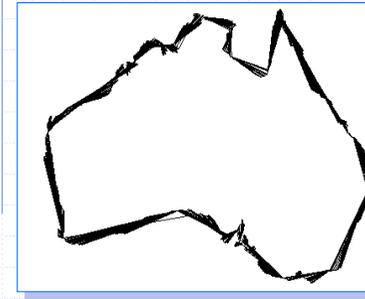
Iterative reduced search for multi-object min- ε problem

Step 1: Find preliminary approximation of every object for an initial number of segments.

Step 2: Iterate the following:

- Apply multiple-goal reduced search DP to define the Rate-Distortion functions.
- Solve the optimal allocation of the segments number among the objects using the Rate-Distortion functions.

Min-# solution as shortest path in digraph

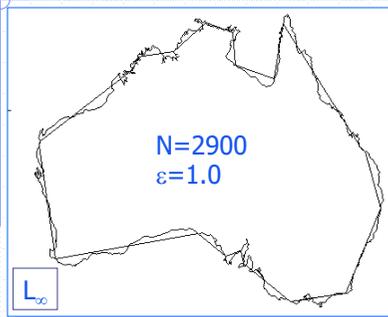


Feasibility graph $G_\varepsilon(P)$
for $\varepsilon=10$ with L_∞ measure.

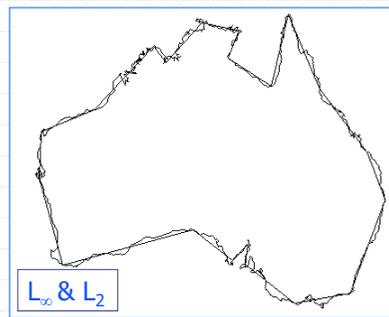


Min-# solution as shortest
path in the digraph $G_\varepsilon(P)$.

Joint using of L_∞ & L_2 error measures



A shortest path in digraph $G_\varepsilon(P)$
for $\varepsilon=1.0$: $E_2=629$, $d_{\max}=1.0$;
 $T=4.5$ s.



After optimization with reduced
search: $E_2=298$, $d_{\max}=1.08$;
 $T=5.0$ s.

Main results

- Iterative Reduced Search algorithm, including:
 - Algorithm for min- ε problem
 - Algorithm for multiple-object min- ε problem
 - Algorithm for min-# problem with L_∞ & L_2 error measures
- Approximation algorithm for closed curves with state space analysis

Email: koles@cs.joensuu.fi

URL: http://cs.joensuu.fi/~koles/approximation/Ch3_0.html