

# TCP, HTTP and Java over Wireless Links

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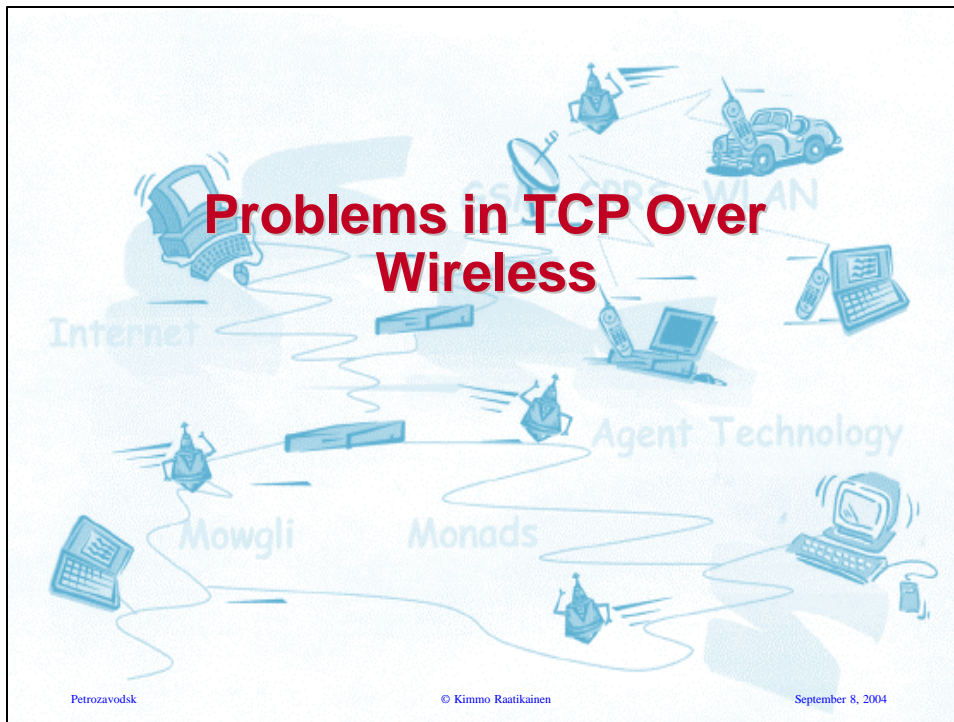
## Lessons Outline

- Problems in TCP Over Wireless
- Additional Problems Due to HTTP
- Java RMI Over Wireless

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2



**Problems in TCP/IP over Wireless Link - 1**

- Overhead due to protocol headers
  - TCP headers take c. 60 bytes
- High latency
  - “extra” round trips should be avoided
- TCP slow start
  - full bandwidth not utilized
- Timers will not work as intended
  - If packet delivery times vary, then TCP timers get confused

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## Problems in TCP/IP over Wireless Link - 2

- Inefficient recovery from packet losses
- Simultaneous TCP connection interfere with each other
- No support for disconnected state

## Should TCP be used – facts against

- It is generally recognized that TCP does not perform well in the presence of significant levels of non-congestion loss.
  - TCP detractors argue that the wireless medium is one such case, and that it is hard enough to fix TCP. They argue that it is easier to start from scratch.
- TCP has too much header overhead.
- By the time the mechanisms are in place to fix it, TCP is very heavy, and ill-suited for use by lightweight, portable devices.

## Should TCP be used – facts in favor – 1/2

- It is preferable to continue using the same protocol that the rest of the Internet uses for compatibility reasons.
  - Any extensions specific to the wireless link may be negotiated.
- Legacy mechanisms may be reused (for example three-way handshake).
- Link-layer FEC and ARQ can reduce the BER such that any losses TCP does see are, in fact, caused by congestion (or a sustained interruption of link connectivity).
  - Modern W-WAN technologies do this (CDPD, US-TDMA, CDMA, GSM), thus improving TCP throughput.

## Should TCP be used – facts in favor – 2/2

- Given TCP's wealth of research and experience, alternative protocols are relatively immature, and the full implications of their widespread deployment not clearly understood.
- Handoffs among different technologies are made possible by Mobile IP [RFC2002], but only if the same protocols, namely TCP/IP, are used throughout.

## What to Improve - 1/2

- TCP: Current Mechanisms
  - Slow Start and Congestion Avoidance
  - Fast Retransmit and Fast Recovery
- Connection Setup with T/TCP [RFC1397, RFC1644]
- Slow Start Proposals
  - Larger Initial Window
  - Growing the Window during Slow Start
    - ACK Counting
    - ACK-every-segment
  - Terminating Slow Start
  - Generating ACKs during Slow Start
- ACK Spacing

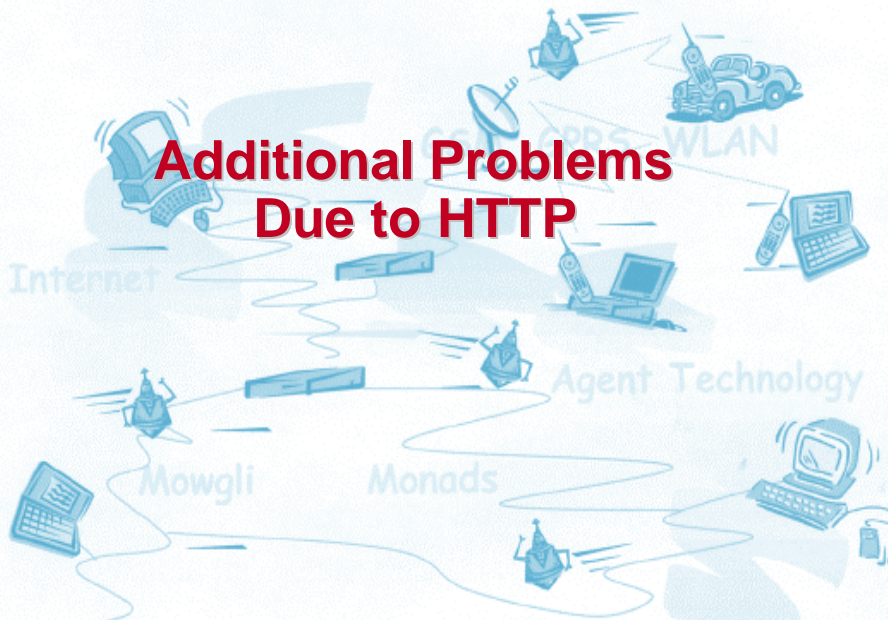
## What to Improve - 2/2

- Delayed Duplicate Acknowledgements
- Selective Acknowledgements [RFC2018]
- Detecting Corruption Loss
  - Without Explicit Notification
  - With Explicit Notifications
- Active Queue Management
- Scheduling Algorithms
- Split TCP and Performance-Enhancing Proxies (PEPs)
- Header Compression Alternatives
- Payload Compression
- TCP Control Block Interdependence

## Future Readings

- RFC 2757: Long Thin Networks
- Other IETF PILC WG RFCs and Ids
- Particularly, TCP over Second (2.5G) and Third (3G) Generation Wireless Networks
  - draft-ietf-pilc-2.5g3g-10

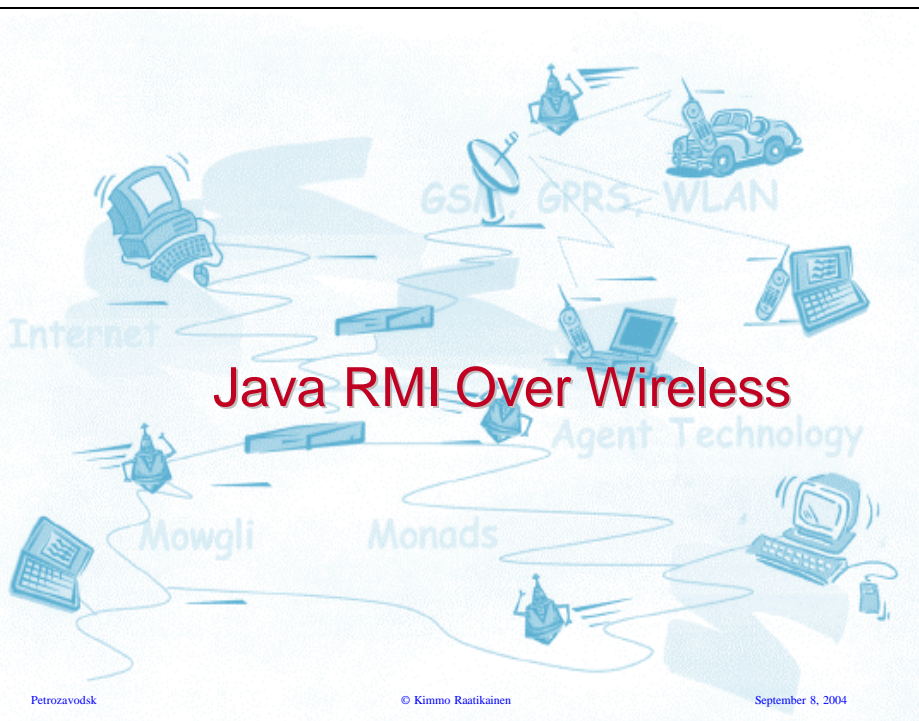
## Additional Problems Due to HTTP



## HTTP Performance Problems

- HTTP requires an excessive number of round-trips
  - In HTTP/1.0 each hypertext document involves creation and deletion of several TCP connections
  - one for each embedded document object
- Connections are typically short
  - large overhead
  - three-way handshake
  - TCP slow-start

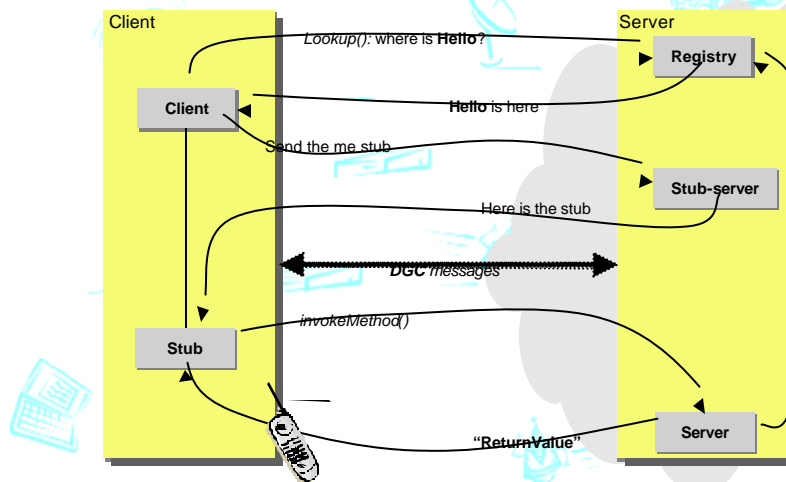
## Java RMI Over Wireless



## Remote Method Invocation

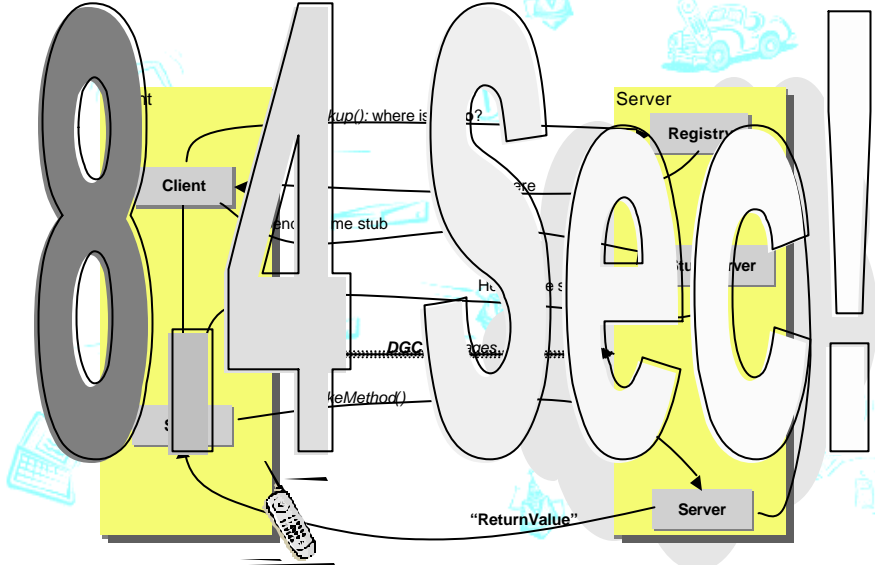
- RMI protocol interface lets Java objects on different hosts communicate with each other in a transparent way
- Clients can invoke methods of a remote object as if they were local methods
- Preserve the object oriented paradigm in distributed computing

## Java RMI in a Nutshell





# Java RMI in a Nutshell

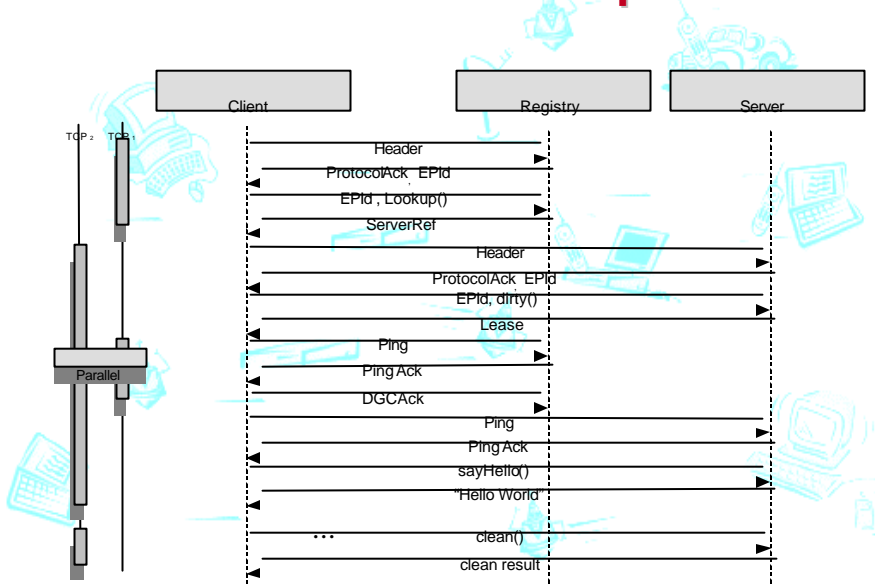


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17

# "Hello World" Example



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18

## Data traffic analysis

	Client to Server and Registry (bytes)	Server and Registry to Client (bytes)	Total (bytes)
Registry Lookup	55 (6%)	276 (42%)	331 (20%)
Invocation Data	41 (4%)	37 (6%)	78 (5%)
DGC Data	831 (85%)	305 (46%)	1136 (69%)
Protocol Overhead	52 (5%)	40 (6%)	92 (6%)
Total	979 (100%)	658 (100%)	1637 (100%)

## RMI Optimization

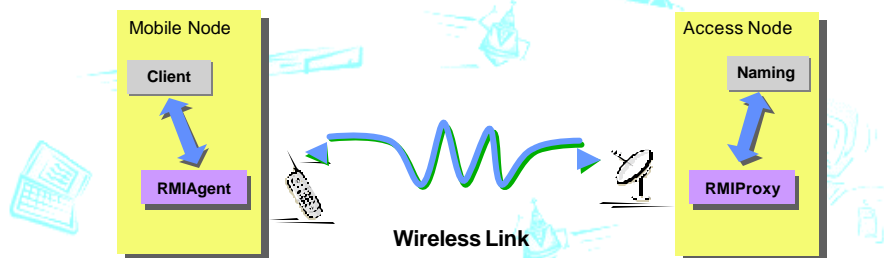
- Maintain compatibility with Java RMI specifications
- Avoid redundancy in communication protocol
- Use compression and caching to minimize data transmission

## Java RMI Optimization

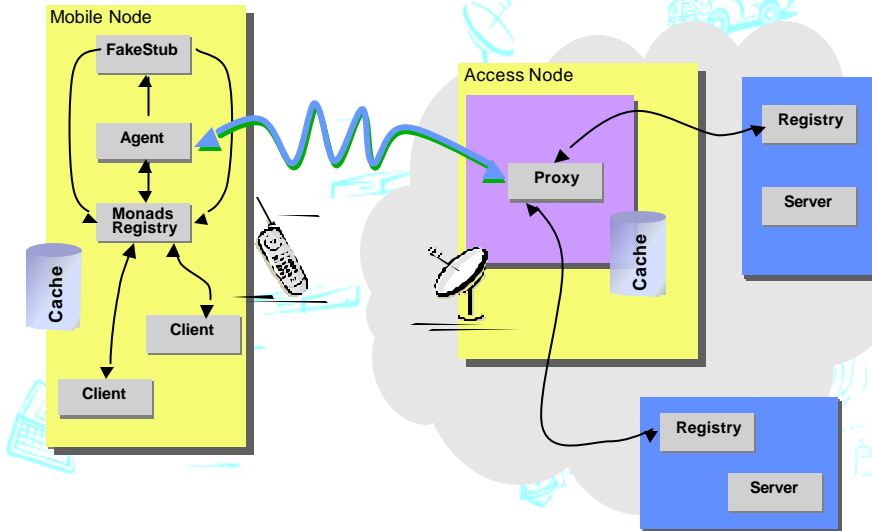
- Protocol
  - Use of Mediators to minimize the exchange of data through the wireless link.
- Data Communication
  - Optimized Communication: Compress and Optimize data communication
- Stub&Class Loading
  - If possible, avoid to download stubs

## Protocol Optimization

- The idea is to de-couple the connection between the client and the server using mediators.



# Optimized RMI

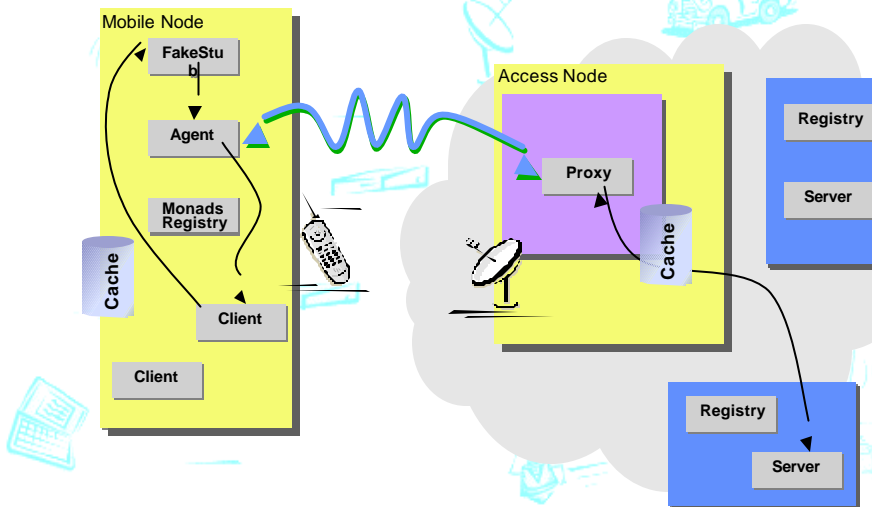


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23

# Optimized RMI

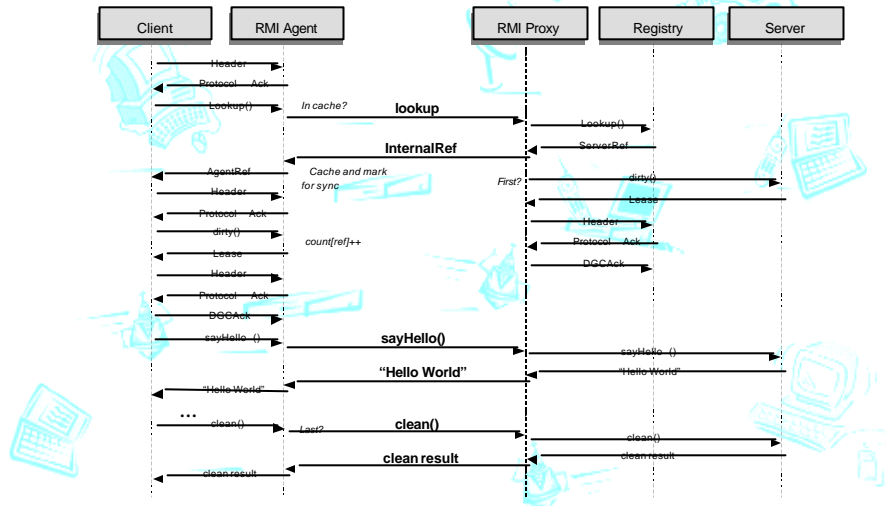


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24

## Optimized Remote Invocation



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## Comparison between Normal RMI and Optimized RMI

	Registry Invocation	Remote Invocation	Total
Java RMI	7.1 sec	1.3 sec	8.4 sec
Optimized RMI	1.7 sec	0.6 sec	2.3 sec
Improvement	417%	216%	365%

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26

## Conclusions

- Designers cannot just “plug-in” wireless communication to existing solutions
- Wireless issues extend their influence also to middleware component and eventually to applications
- Solutions are there, just mostly ignored

## Wireless Java RMI Background



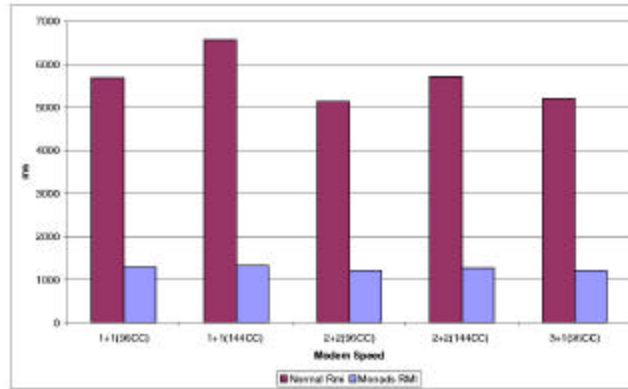
## Test Arrangements

- Operating Systems
  - Clients:
    - Windows98
    - Linux (Red Hat 6.1, kernel 2.2.14)
  - Server
    - Windows NT (SP 6)
    - Linux (Red Hat 6.1, kernel 2.2.14)

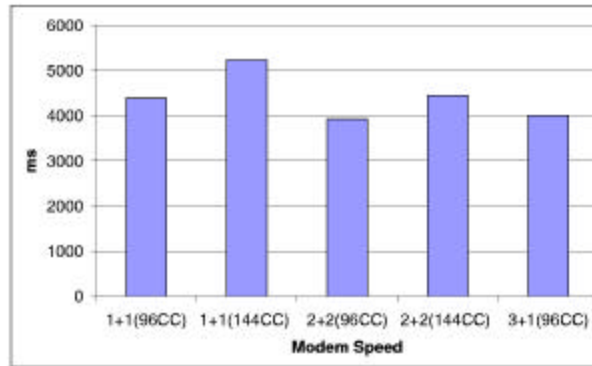
## Test Arrangements

- Java Virtual Machine
  - Sun JDK 1.2.2 (Linux and Windows)
- Wireless communication
  - GSM HSCSD (5 configurations)
- Benchmark Suite
  - KaRMI from University of Karlsruhe

## Lookup Results (windows)

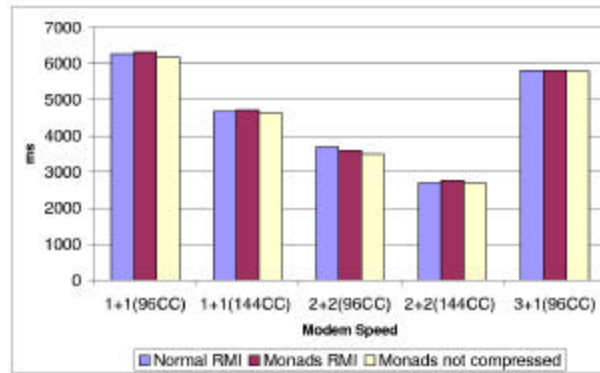


## Lookup Differences





## Invocation Results Image Uplink (Linux)



## Invocation Results Two-way Text uplink

